

Endomorphism Monoids of Combinatorial Structures

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Relational structures

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- ▶ The sequence $(r_i (i \in I))$ is called the **signature** of \mathcal{X} .

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- ▶ Posets: signature (2) ; properties R, AS, T.
- ▶ Permutations: signature $(2,2)$; properties: two linear orders.

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Let $\mathcal{X} = (X; R_i (i \in I))$ be a relational structure. An **endomorphism** is a mapping $\theta : X \rightarrow X$ which respects all the relations R_i , i.e.

$$(x_1, \dots, x_k) \in R_i \Rightarrow (x_1\theta, \dots, x_k\theta) \in R_i.$$

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General Problem

For a given \mathcal{X} , how are \mathcal{X} , $\text{End}(\mathcal{X})$ and $\text{Aut}(\mathcal{X})$, and their properties, related?

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Problem

Describe the monoids $\text{End}(P)$ where P is a permutation.

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Question

What is $\text{rank}(B_n)$, where B_n is the monoid of all binary relations on $\{1, \dots, n\}$?

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Example

$\text{rank}(\text{Trans}(n) : \text{Sym}(n)) = 1$.

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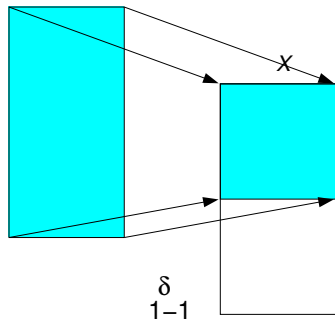
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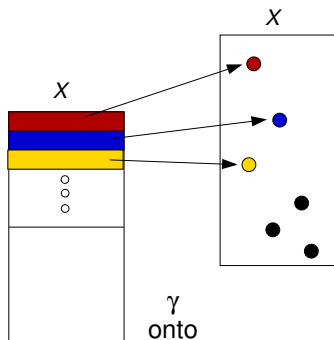


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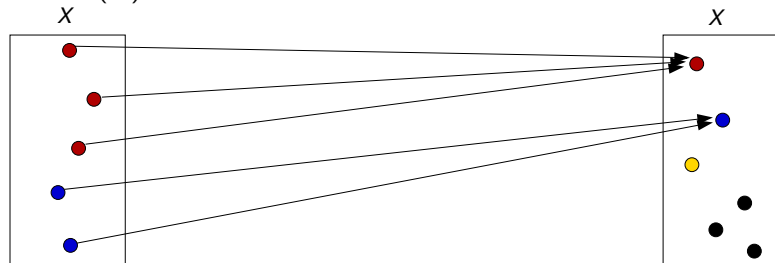


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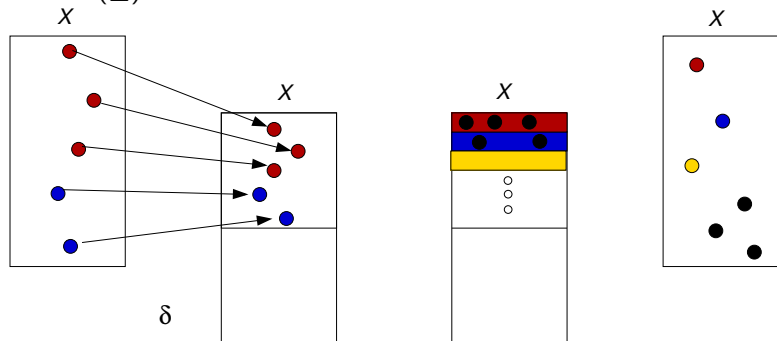


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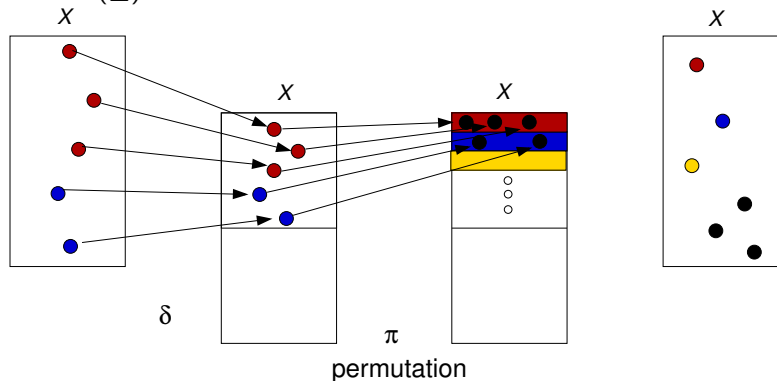


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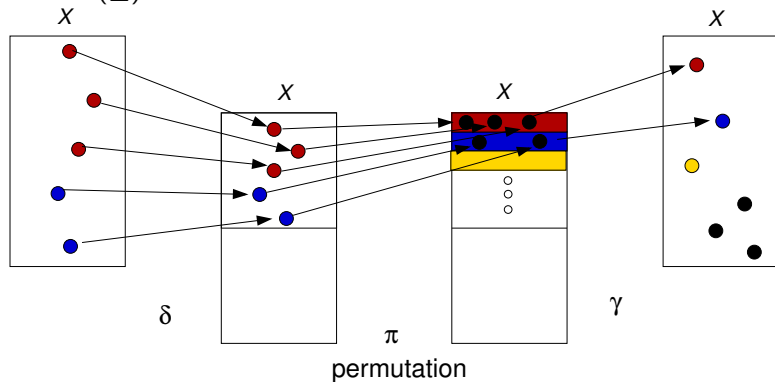


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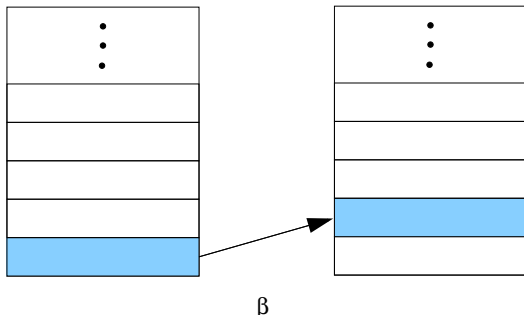
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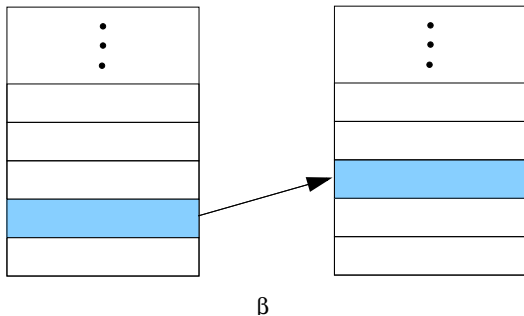


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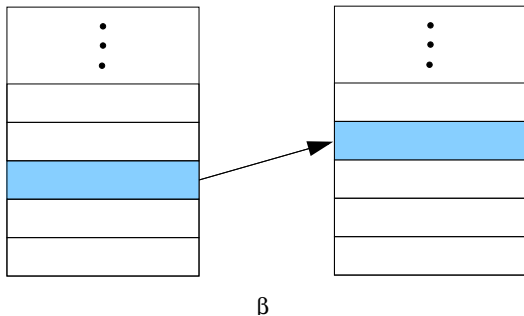


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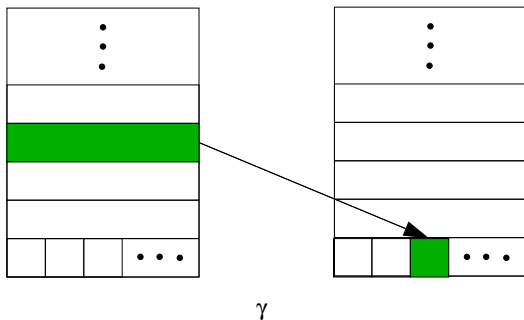


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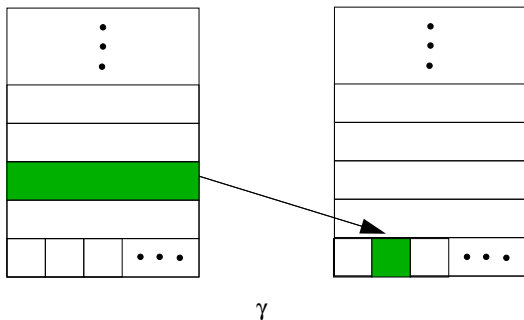


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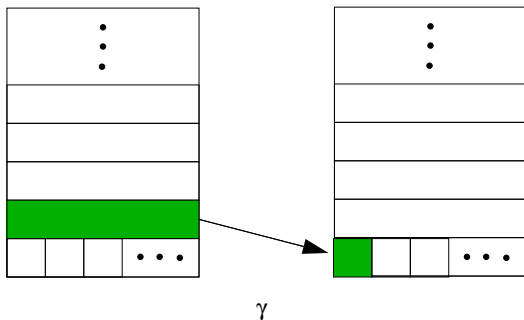


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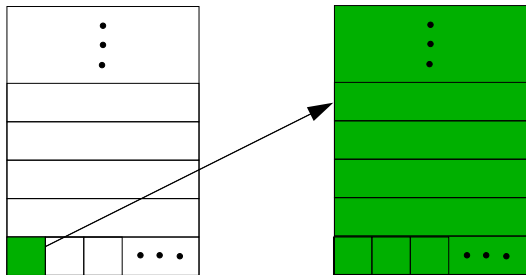


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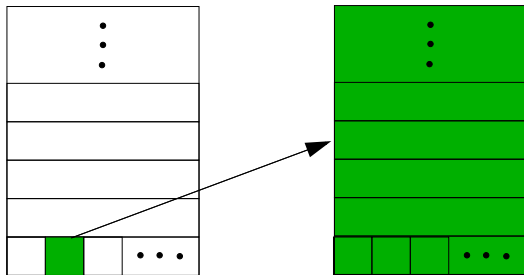
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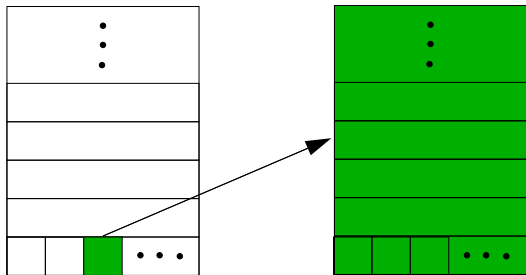
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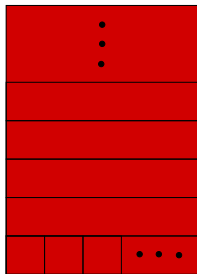
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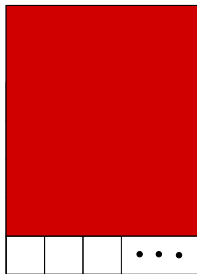
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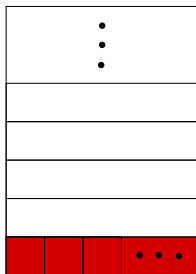
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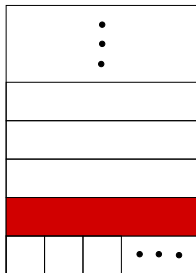
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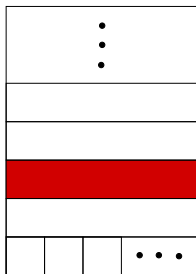
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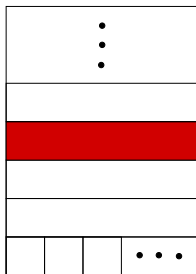
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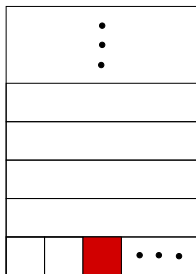
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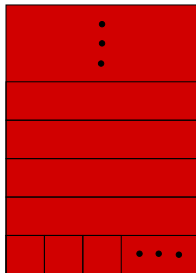
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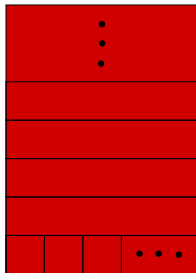
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We say that $T(X)$ has **Sierpinski index 2**.

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Problem

What is the Sierpinski index of $C((0, 1))$?

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Theorem (Higgins, Howie, Mitchell, NR 03)

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Then $\gamma = \epsilon\delta \in \langle \text{End}(\mathbb{N}), \delta \rangle$. □

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Question

Classify, in order-theoretic terms, all linearly ordered sets L with finite relative ranks in $\text{Trans}(L)$.

Partially ordered sets

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Let P be a partially ordered set. For $x \in P$ write:

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- (III) If P is an antichain, then $\text{rank}(\text{Trans}(P), \text{End}(P)) = 0$:-)

Partially ordered sets

Corollary

$\text{rank}(\text{Trans}(P), \text{End}(P)) \leq 2$ provided any of the following hold:

- ▶ P has a smallest element;
- ▶ P has a largest element;
- ▶ P is a lattice.
- ▶ ...

A poset with relative rank 2 (I)

Higgins, Mitchell, Morayne, NR 06

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If P is a poset with no non-trivial mono- or epi-morphisms then $\text{rank}(\text{Trans}(P), \text{End}(P)) \geq 2$.

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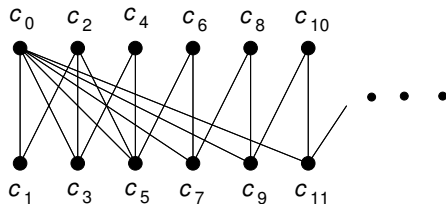
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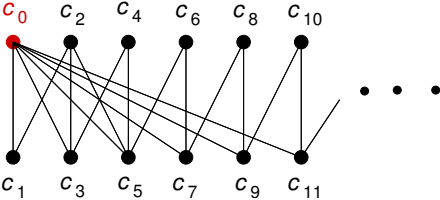
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By induction $\pi = \mu^n$, a contradiction.

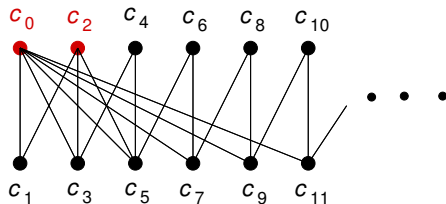
Poset with no monomorphisms



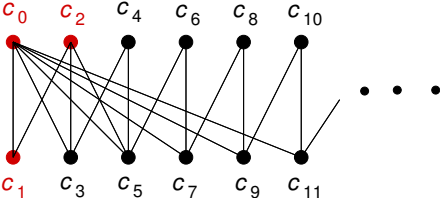
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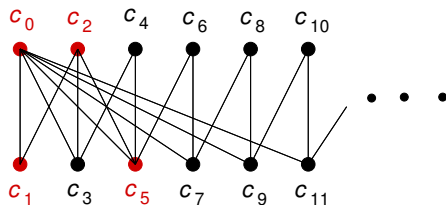
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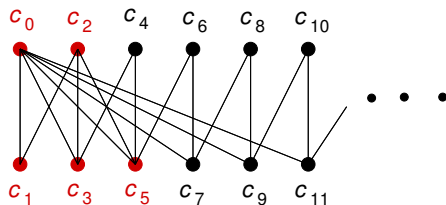
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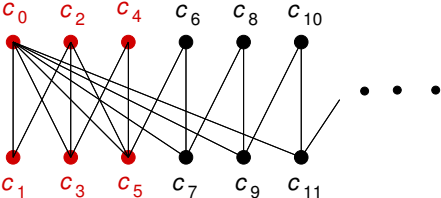
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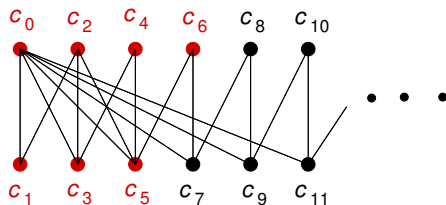
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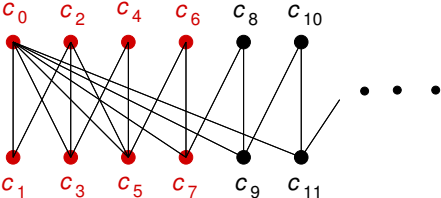
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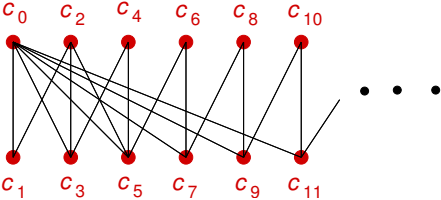
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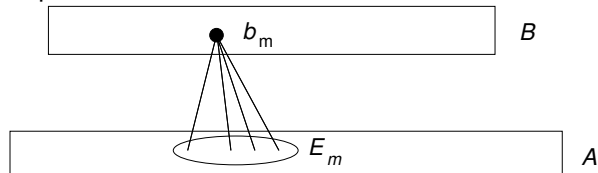
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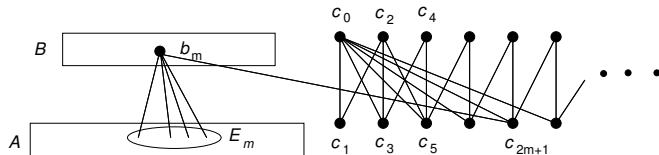
Enumerate \mathcal{E} : E_1, E_2, \dots

Let $B = \{b_1, b_2, \dots\}$.

A partial order on $A \cup B$: $b_m > a_n \Leftrightarrow a_n \in E_m$.



A poset with relative rank 2 (II)



Proposition (Higgins, Mitchell, Morayne, NR 06)

The above poset P satisfies the conditions from the part (I) of the Theorem, but has no non-trivial mono- or epi-morphisms; consequently $\text{rank}(\text{Trans}(P) : \text{End}(P)) = 2$.

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Definition

The **random graph** R is the unique countable graph with the following property:

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R is universal: R contains a copy of every countable graph.

In fact, homogeneity and universality provide another definition of R .

$\text{Aut}(R)$ and $\text{End}(R)$: Similar?

Theorem (Truss 85)

Every countable group embeds into $\text{Aut}(R)$.

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R is *not* characterised by the properties of universality and homomorphism homogeneity.

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To every ideal I of a monoid M there corresponds the Rees congruence $\Phi_I \cup \Delta_M$ on M ; but, not every congruence is a Rees congruence.

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Question

How many non-Rees congruences does End(R) have?

Random graph: Questions

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Does $\text{Aut}(R)$ have a finite Sierpinski index?

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